

DNS security

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Tuesday, September 18, 2012

1 Chain of Trust

- How the Chain of Trust works
- Details of RR's used in the chain of trust

2 Denial of existence

- Concept and Resource Records
- Details of RR's used in denial of existence

3 Resolver and protocol issues

Outline

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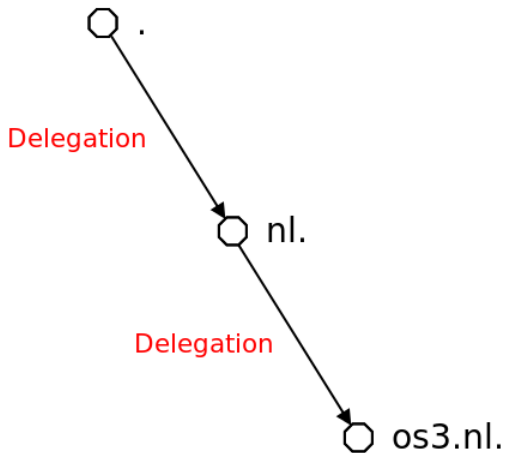
3 Resolver and protocol issues

Tree walking algorithm

Validating a RRset for `www.os3.nl`:

- Start at the root
- Verify authenticity of delegation to `nl.` zone
- Verify authenticity of delegation to `os3.nl.` zone
- Verify authenticity of RRset of `www.os3.nl`

Tree walking algorithm

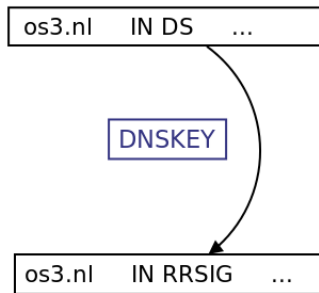
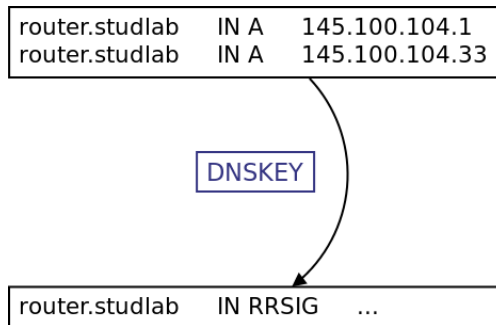


Tree walking algorithm

Resource Records used to build the Chain of Trust:

- RRSIG
 - Resource Record Signature
 - Contains signatures of RRsets
- DNSKEY
 - Stores public key of a zone
 - Used to verify signatures
- DS
 - Delegation Signer
 - Contains a hash of the DNSKEY of a child zone

Tree walking algorithm



Tree walking algorithm

Validating delegation from root to nl. zone:

- 1 Keys of root zone are known to resolver
- 2 Ask root zone for DS of nl. zone
- 3 Ask nl. zone for DNSKEY of nl. zone
- 4 Verify that DS is valid hash of DNSKEY

Tree walking algorithm

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Two steps are missing...

Tree walking algorithm

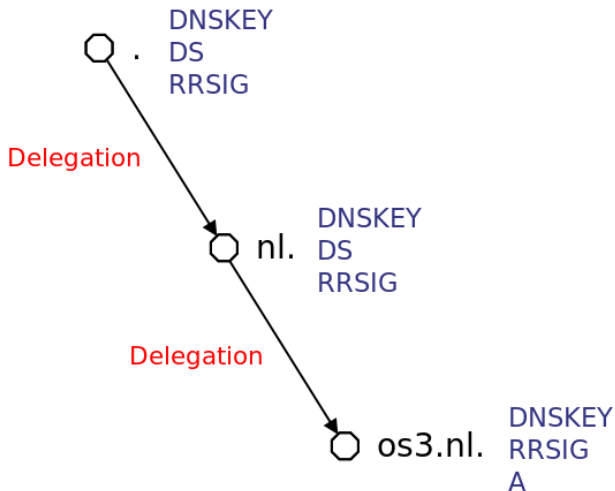
Validating delegation from root to nl. zone:

- 1 Keys of root zone are known to resolver
- 2 Ask root zone for DS of nl. zone
- 3 Ask nl. zone for DNSKEY of nl. zone
- 4 Verify that DS contains a valid hash of DNSKEY

Two steps are missing...

- 5 Retrieve RRSIG for root zone, and
- 6 verify authenticity of DS

Tree walking algorithm



Tree walking algorithm

Validating delegation from root to nl. zone:

- 1 Retrieve DNSKEY for root zone
 - Stored in resolver as secure starting point
 - Can also be retrieved from root nameservers:
 - `dig . DNSKEY`
- 2 Ask root zone for DS of nl. zone
 - Contains a hash of the DNSKEY of nl. zone
 - Stored in root zone
 - `dig nl DS`

Tree walking algorithm

Validating delegation from root to nl. zone:

- 3 Ask nl. zone for DNSKEY of nl. zone
 - `dig nl DNSKEY`
- 4 Verify that DS contains a valid hash of DNSKEY
 - Create hash of DNSKEY, and
 - compare to hash in DS record

Tree walking algorithm

Validating delegation from root to nl. zone:

- 5 Retrieve RRSIG for root zone
 - Stored in root zone
 - `dig . RRSIG`
- 6 Use DNSKEY to verify signature in RRSIG matches data in DS record

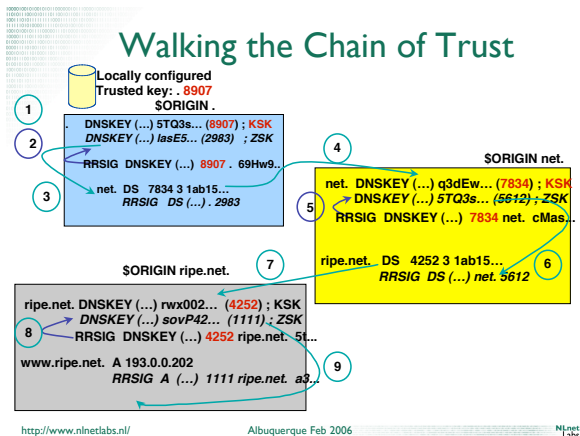
DNSSEC-bis chain of trust (1)

- To validate a resource record set RRset
 - Validate RRSIG(RRset)
 - by using a ZSK (zone signing key) from the DNSKEY RRset
 - Validate RRSIG(DNSKEYset)
 - by using a KSK (key signing key) from that same DNSKEY RRset
 - Validate the KSK
 - by using a DS (present in the parent zone)
which contains a hash of the KSK
 - the KSK used is called a SEP (Secure Entry Point)

DNSSEC-bis chain of trust (2)

- Continue validating one level higher in the hierarchy
 - Use DS as RRset and iterate
- Use trusted anchors for DNSKEY's or DS's to terminate
 - for instance for checking root zone keys
- Authority around a cut is now as follows
 - NS and DNSKEY are authoritative on the **child** side of the cut (zone apex)
 - DS is authoritative on the **parent** side of the cut (delegation point)

Chain of Trust illustration



(Source: Olaf Kolkman, RIPE NCC, NLnet Labs)

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RRSIG record example from RFC 4034

```
host.example.com. 86400 IN RRSIG A 5 3 86400 20030322173103 (
    20030220173103 2642 example.com.
    oJB1W6WNGv+ldvQ3WDGOMQkg5IEhjRip8WTr
    PYGv07h108dUKGMeDPKijVCHX3DDKdfb+v6o
    B9wfuh3DTJXUAFI/M0zm0/zz8bW0Rzn1803t
    GNazPwQKkRN20XPXV6nwwf oXmJQbsLNrLfkG
    J5D6fwFm8nN+6pBzeDQfsS3Ap3o= )
```

RRSIG record content

| Part | Meaning |
|--------------------------|---|
| Type covered | Record type this RRSIG is about |
| Algorithm | Signature algorithm used; 5 is RSA/SHA-1 |
| Labels | Number of labels of owner (without root) |
| TTL | Original time to live |
| Expiration and Inception | Signature validity date bounds |
| Key Tag | To help find (not identify) the signing key |
| Signer's Name | Owner name of zone and key to use |
| Signature | Signature (in Base64) |

DNSKEY record example from RFC 4034

```
example.com. 86400 IN DNSKEY 256 3 5 ( AQPskMynfzW4kyBv015MUG2DeIQ3
      Cbl+BBZH4b/OPY1kxkmvHjcZc8no
      kfzj31GajIQKY+5CptLr3buXA10h
      WqTkF7H6Rf oRqXQeogmMHfpftf6z
      Mv1LyBUgia7za6ZEz0JB0zt yvhjL
      742iU/TpPSEDhm2SNKLi jfUppn1U
      aNvv4w== )
```

DNSKEY record content

| Part | Meaning |
|------------|---|
| Flags | Zone key (KSK+ZSK); Secure entry point (KSK) |
| Protocol | Always 3 (for backward compatibility with KEY RR) |
| Algorithm | Signature algorithm used; 5 is RSA/SHA-1 |
| Public key | Key used for signing (in Base64) |

DS record example from RFC 4034 and 4509

```
dskey.example.com. 86400 IN DNSKEY 256 3 5 ( AQ0eiiROGOMYkDshWoSKz9Xz
fwJr1AYtsmx3TGkJanXVbfi/
2pHm822aJ5iI9BMzNXxeYcmZ
DRD99WYwYqUSdjMmmAphXdvx
egXd/M5+X70rzKBaMbCVdFLU
Uh6DhweJBjEVv5f2wwjM9Xzc
nOf+EPbtG9DMbMADjFDc2w/r
ljwvFw==
) ; key id = 60485
```

```
dskey.example.com. 86400 IN DS 60485 5 1 (
2BB183AF5F22588179A53B0A98631FAD1A292118 ) ; SHA-1
```

```
dskey.example.com. 86400 IN DS 60485 5 2 (
D4B7D520E7BB5F0F67674A0C
CEB1E3E0614B93C4F9E99B83
83F6A1E4469DA50A ) ; SHA-256
```


DS record content

| Part | Meaning |
|---------------|--|
| Key Tag | To help find (not identify) the signing key |
| Algorithm | Signature algorithm of the signing key |
| Digest Type | Hashing algorithm used; 1 is SHA-1, 2 is SHA-256 |
| Digest | Sequence of case-insensitive hexadecimal digits |

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Positive and negative answers

- Proving existence (and authenticity) is elaborate, but “easy”
 - It uses standard public key cryptography signing mechanisms
- Proving non-existence is much more difficult
 - There is nothing to sign
 - Empty non-terminals
 - Wildcards

Methods for denial of existence

- Doing a zone transfer
 - Only possible for trusted clients
- On the fly (realtime) signing of negative replies
 - Often too compute intensive
- Enumerating the zone
 - Makes “zone walking” possible
 - Unless the real names are hidden somehow

More DNSSEC-bis resource records

- **NSEC**: authenticated denial of existence
 - Uses canonical ordering to list domain names in a zone
 - Lexicographic ordering for labels by considering them as “unsigned left-justified octet strings”
 - Lexicographic ordering for domain names as sequence of labels, in fact walking the domain name tree in preorder traversal
- **NSEC3** introduced later to frustrate zone enumeration
 - All domain names are disguised as hashes
 - One can still count the number of domain names

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NSEC record example from RFC 4034

```
alfa.example.com. 86400 IN NSEC host.example.com. (  
    A MX RRSIG NSEC TYPE1234 )
```


NSEC record content

| Part | Meaning |
|------------------|------------------------------------|
| Next Domain Name | Defined by a standard ordering |
| Type Bit Maps | Types present at the current owner |

- This mechanism can be used to enumerate a complete zone...
 - ...which is why certain organisations refuse to implement DNSSEC-bis
- NSEC chains wrap around and fork or join at a delegation point or zone apex
- NSEC3 RR's can anonymise the NSEC records

NSEC3 record example from RFC 5155

```
q04jkcevqvmu85r014c7dkba38o0ji5r.example. NSEC3 1 1 12 aabbccdd (  
    r53bq7cc2uvmubfu5ocmm6pers9tk9en A RRSIG )
```

NSEC3 record content

| Part | Meaning |
|------------------------|---|
| Hash Algorithm | 1 for SHA-1 |
| Flags | Opt-out for skipping unsigned delegations |
| Iterations | Number of times hash is calculated |
| Salt Length | 0-255 (not represented in text format) |
| Salt ¹ | Prevents dictionary attacks |
| Hash Length | 1-255 (not represented in text format) |
| Next Hashed Owner Name | (binary data) |
| Type Bit Maps | Types present at the current owner |

¹Hash Algorithm, Flags, Iterations and Salt are communicated to slave servers via the NSEC3PARAM Resource Record in the zone apex.

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Protocol flags

- DO (DNSSEC OK)
 - Indicate that you want all DNSSEC-related information like RRSIG and NSEC(3) for normal queries, if available
 - Presupposes use of EDNS0
- CD (Checking Disabled)
 - From resolver to (recursive) name server
 - Indicate that you want to do checking yourself
 - Also insecure (invalid) data is passed along
- AD (Authentic Data)
 - From recursive name server to resolver
 - Data has been verified successfully
 - In most cases not set by authoritative server
 - Checking is done at the resolver side

Security status of RR's

- **Secure**
 - Chain of trust exists and signature is valid
- **Insecure**
 - Chain of trust exists and signature is invalid
 - Results in SERVFAIL
- **Bogus**
 - Chain of trust should exist but is broken
 - Results in SERVFAIL
- **Indeterminate**
 - Chain of trust does not exist