

The effects of
disruption
and
topology
on
the computer network
packet delivery

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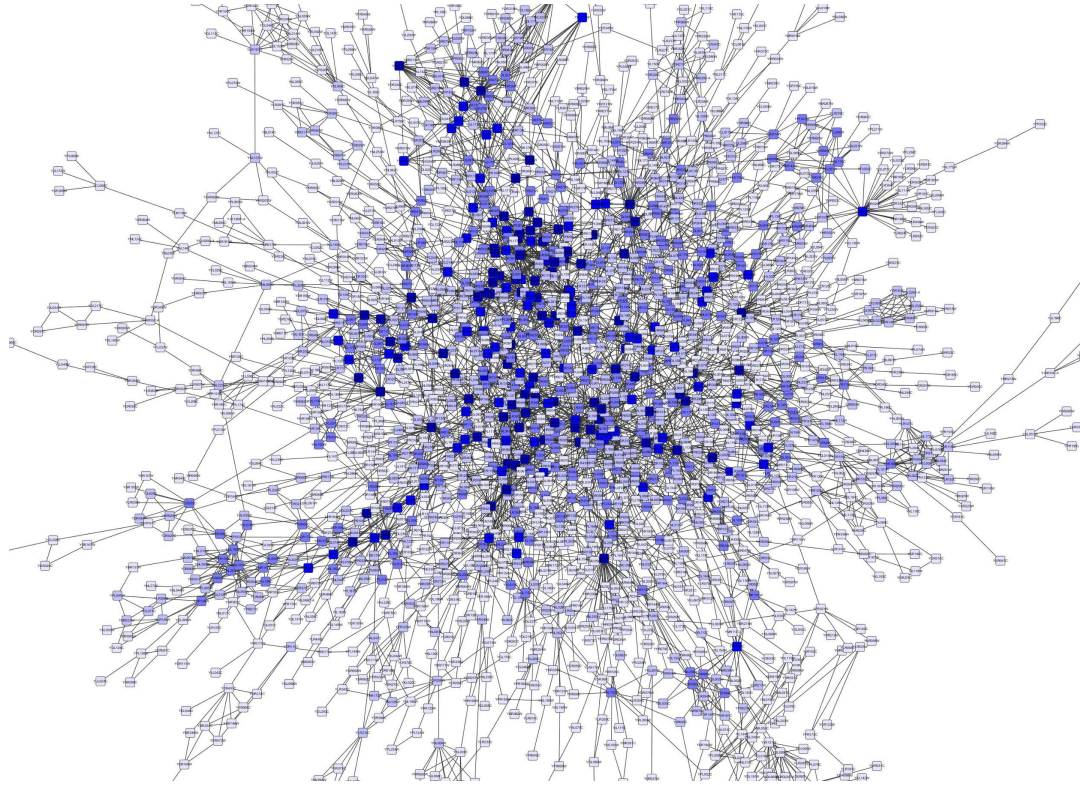
Supervisor : Marc X. Makkes

Networks

Active field of research.

Finding the dependencies between :

- Network “noise”
- Topology
- Transport efficiency

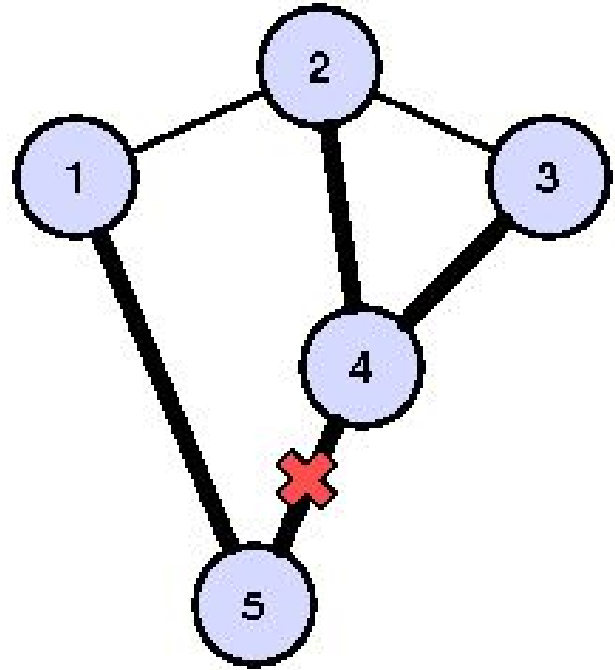


0.6-1.4%

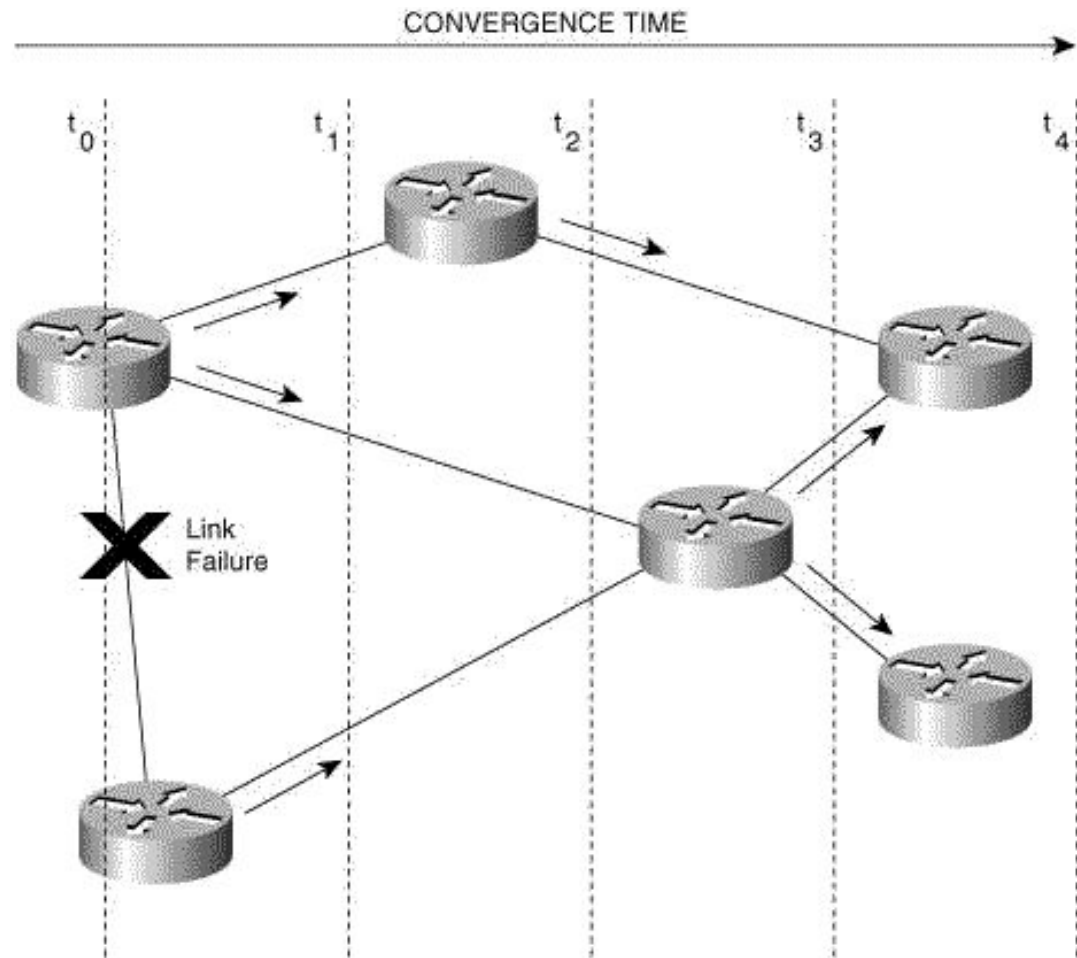
Nygren, Erik, Ramesh K. Sitaraman, and Jennifer Sun. "The Akamai network: a platform for high-performance internet applications." ACM SIGOPS Operating Systems Review 44.3 (2010): 2-19.

Network disruption

- Detect
- Inform others



Convergence takes time...



100ms-15min

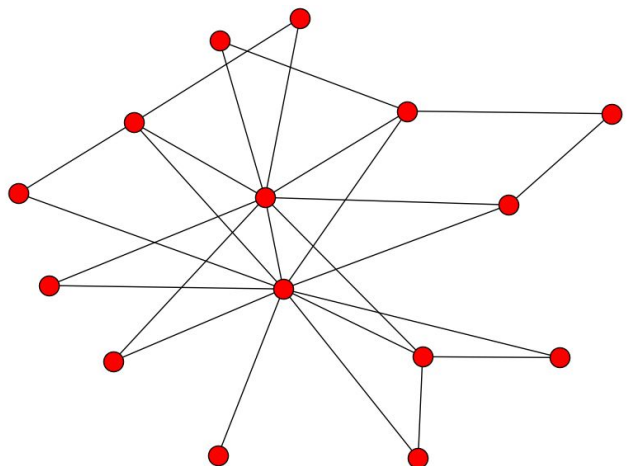
Yan, Hong, et al. "Tesseract: A 4D Network Control Plane." NSDI. Vol. 7. 2007.

Labovitz, Craig et al. "The impact of Internet policy and topology on delayed routing convergence." *INFOCOM 2001. Twentieth Annual Joint Conference of the IEEE Computer and Communications Societies. Proceedings. IEEE 2001*: 537-546.

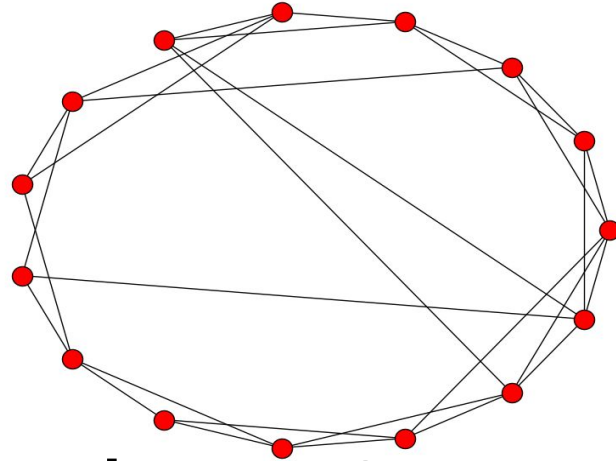
Research questions

How the network resiliency during network reconvergence is affected by :

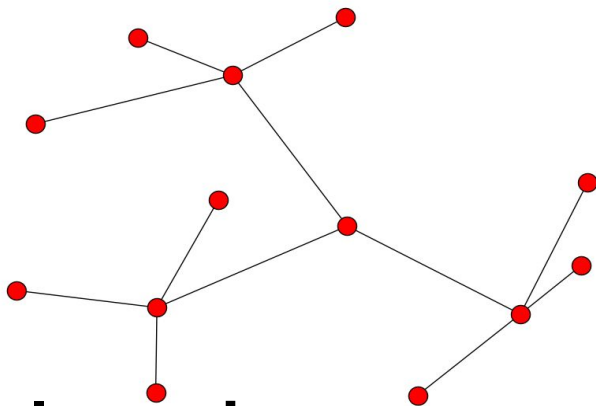
- topology **type** and **size**
- topology **partitioning**
- link state change **probability**



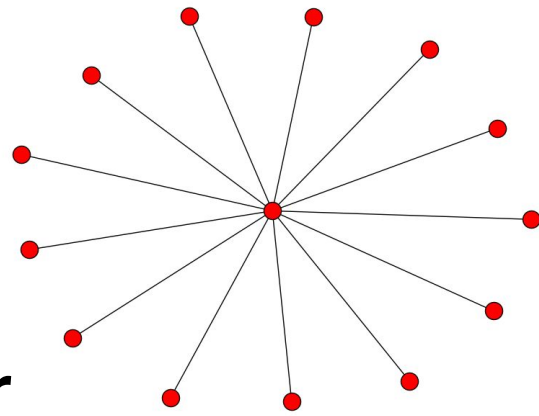
Barabasi-Albert (BA)



Connected Watts-Strogatz (WS)



Balanced tree




Star

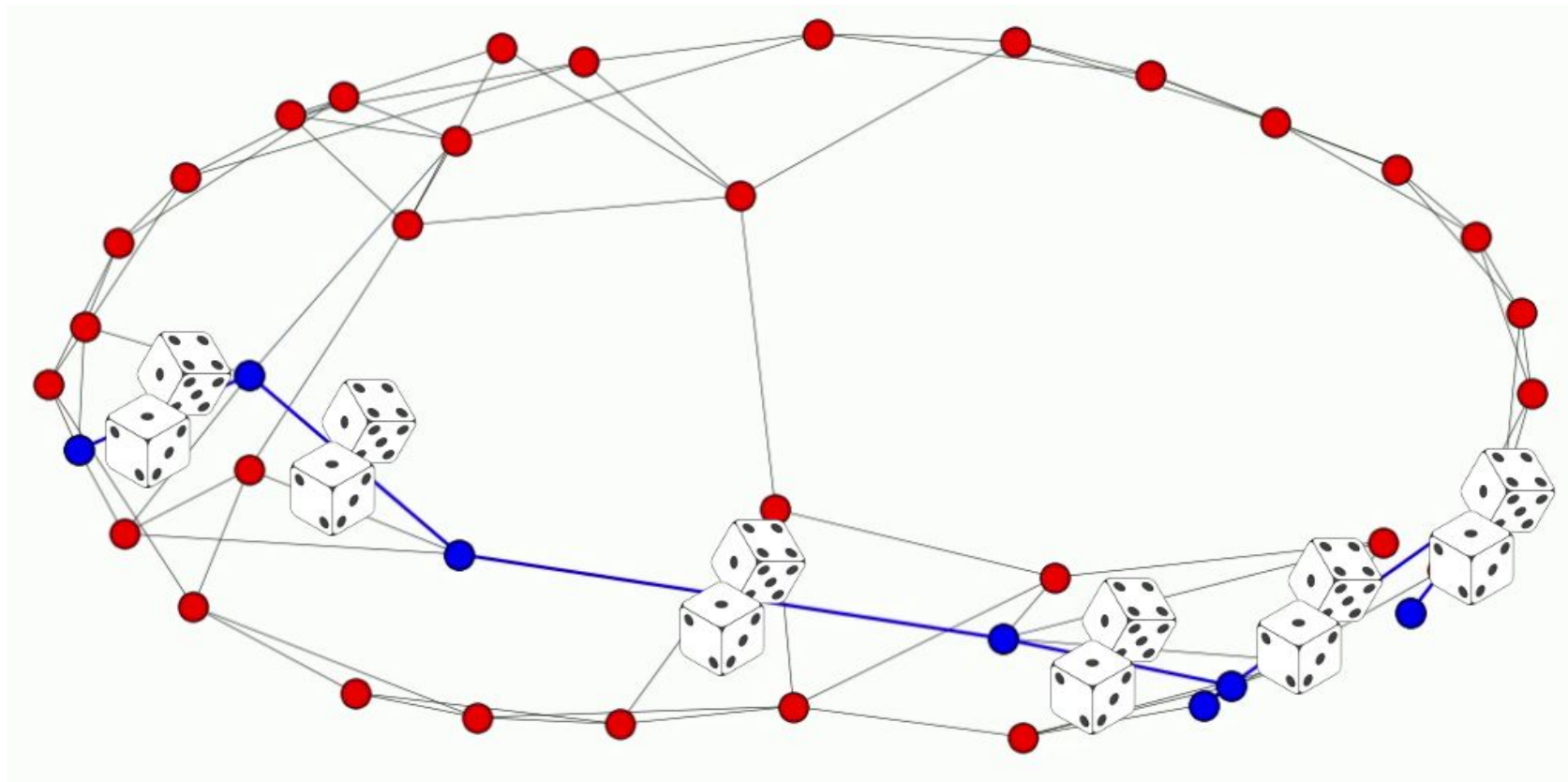
Topology sizes

- Small (40)
- Medium (121)
- Large (364)

... and partitioning

- | | | |
|-----|---|------|
| ● 2 |  | ● BA |
| ● 4 | | ● WS |
| ● 8 | | |

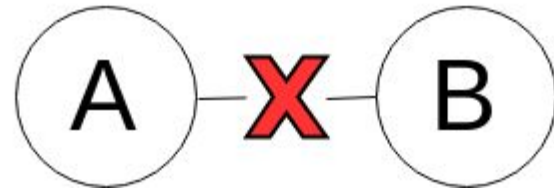
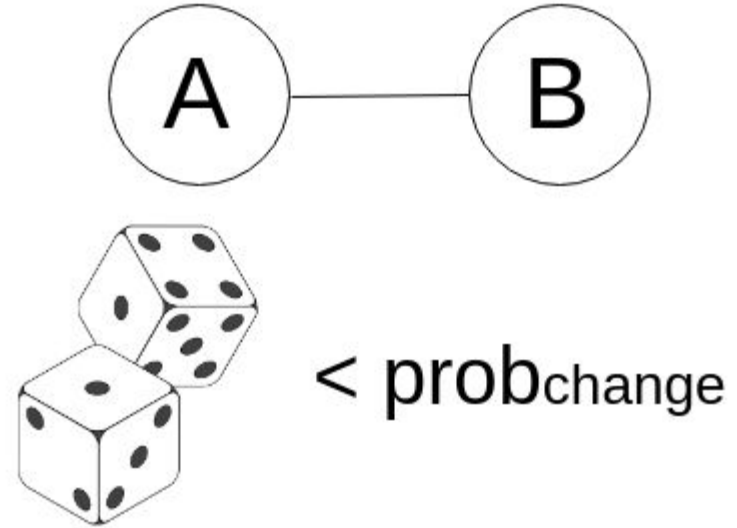
Introducing failures



Probability of link state change

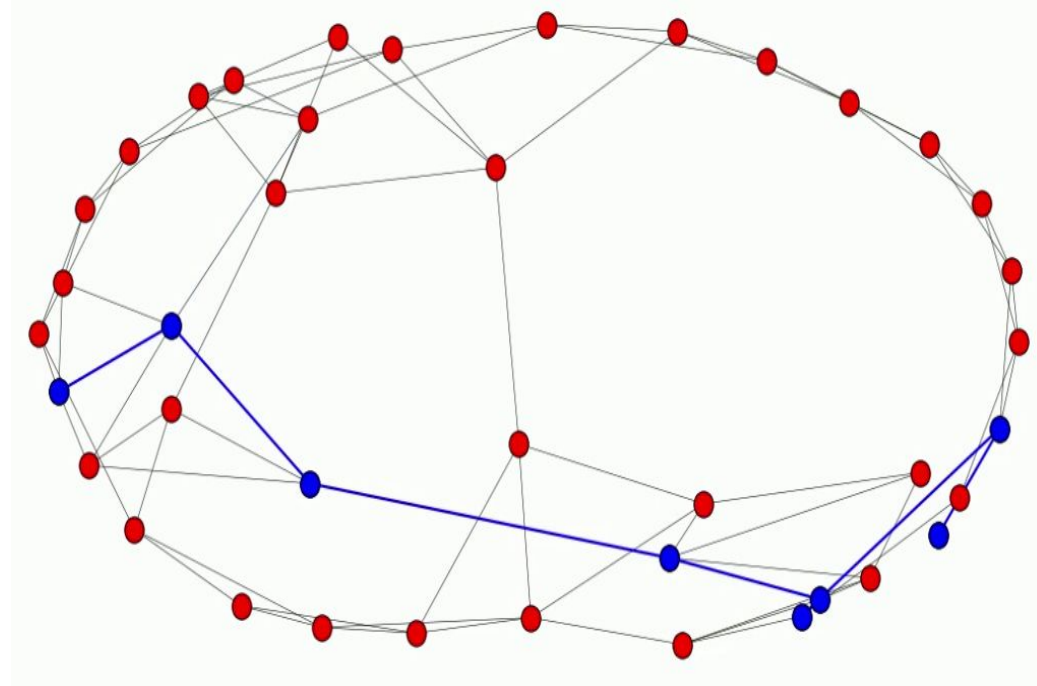
- Experiments conducted for 10 logarithmically spaced values of probability

$\text{prob}_{\text{change}} = \langle 0, 1 \rangle$



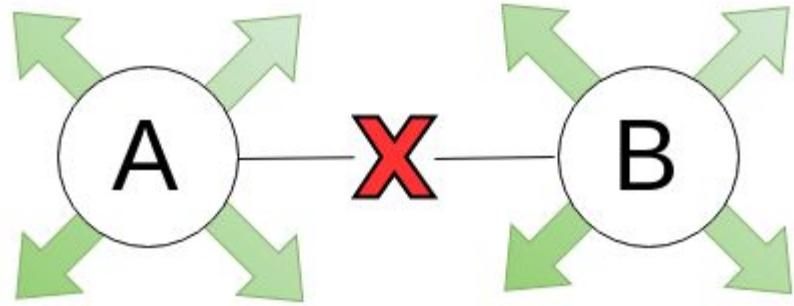
In-simulator link-state protocol

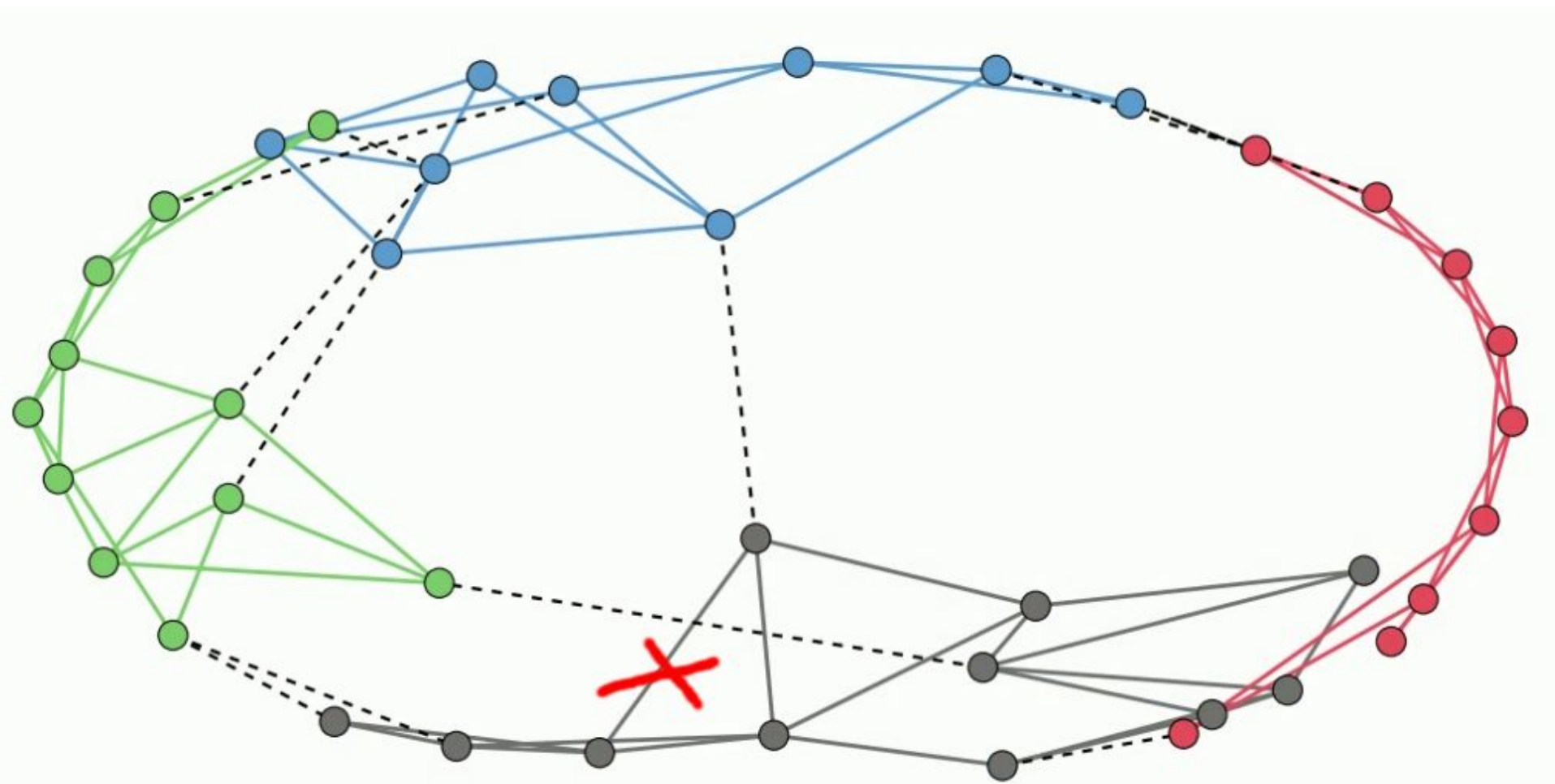
- Every node has knowledge about the **complete** topology
- Runs link-state shortest path routing algorithm

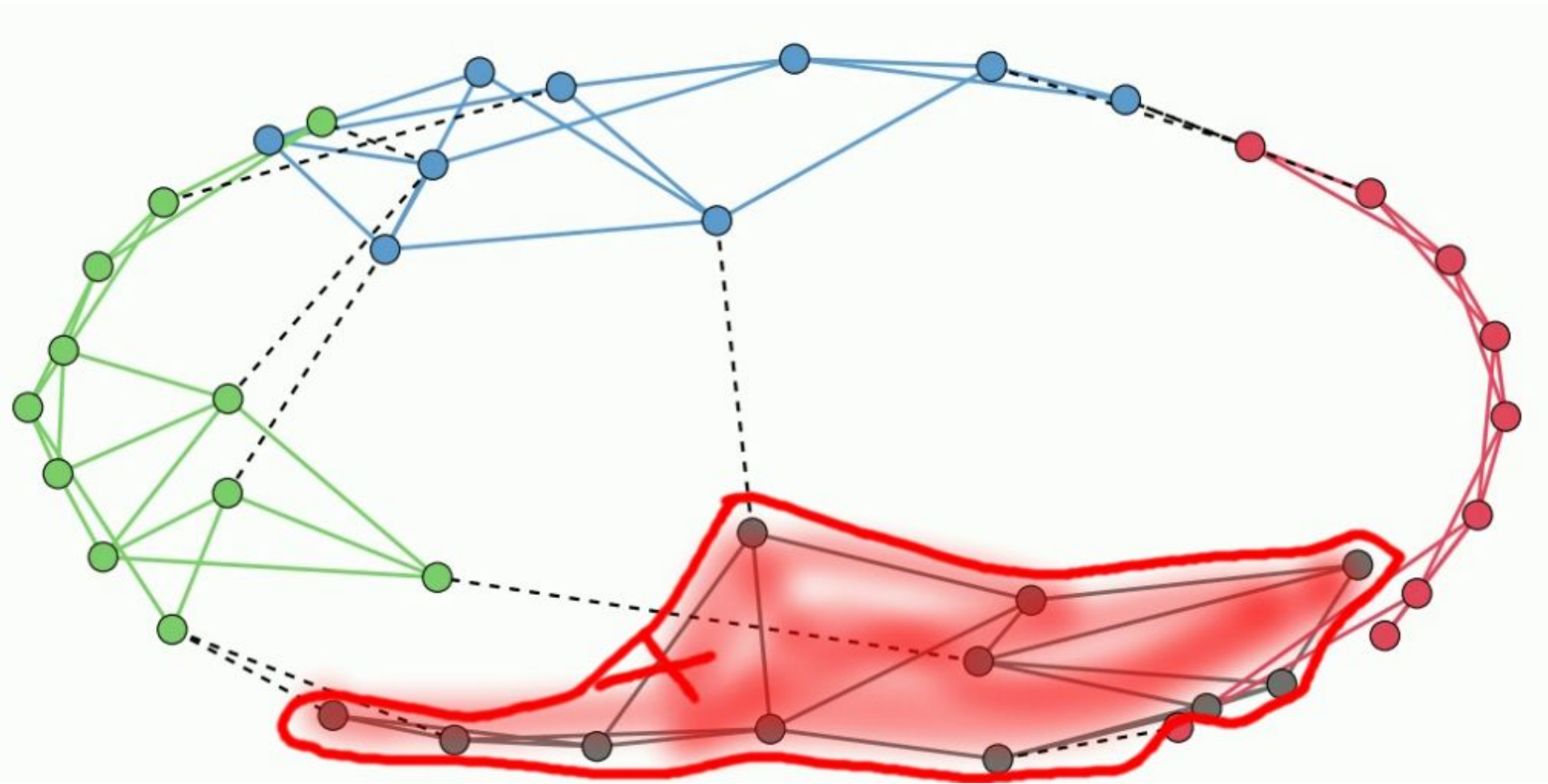


In-simulator link-state protocol

- On the link change event “state” packet sent to the neighbours







Metrics

Delivery ratio

$$\frac{\textit{packets received}}{\textit{packets transmitted}} \times 100 [\%]$$

Mean path length

$$\frac{\sum_{i=0}^N \textit{path length}(i)}{N}$$

$N = \textit{received packet count}$

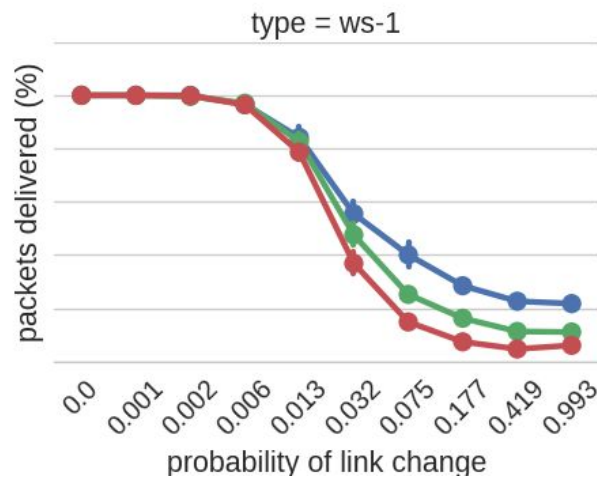
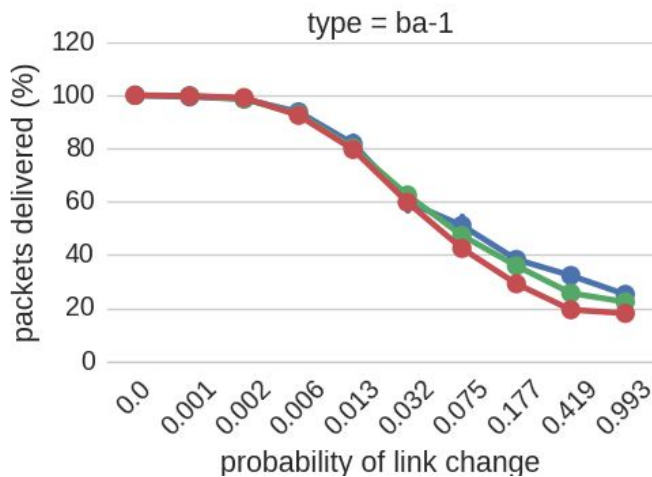
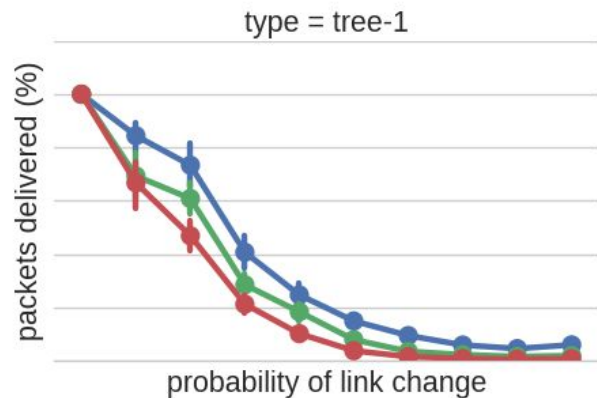
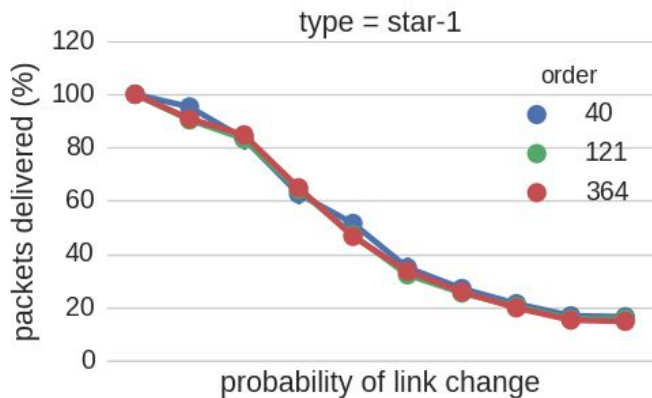
“TTL exceeded” ratio

$$\frac{\textit{TTL dropped packets}}{\textit{packets transmitted}} \times 100 [\%]$$

“NO ROUTE” ratio

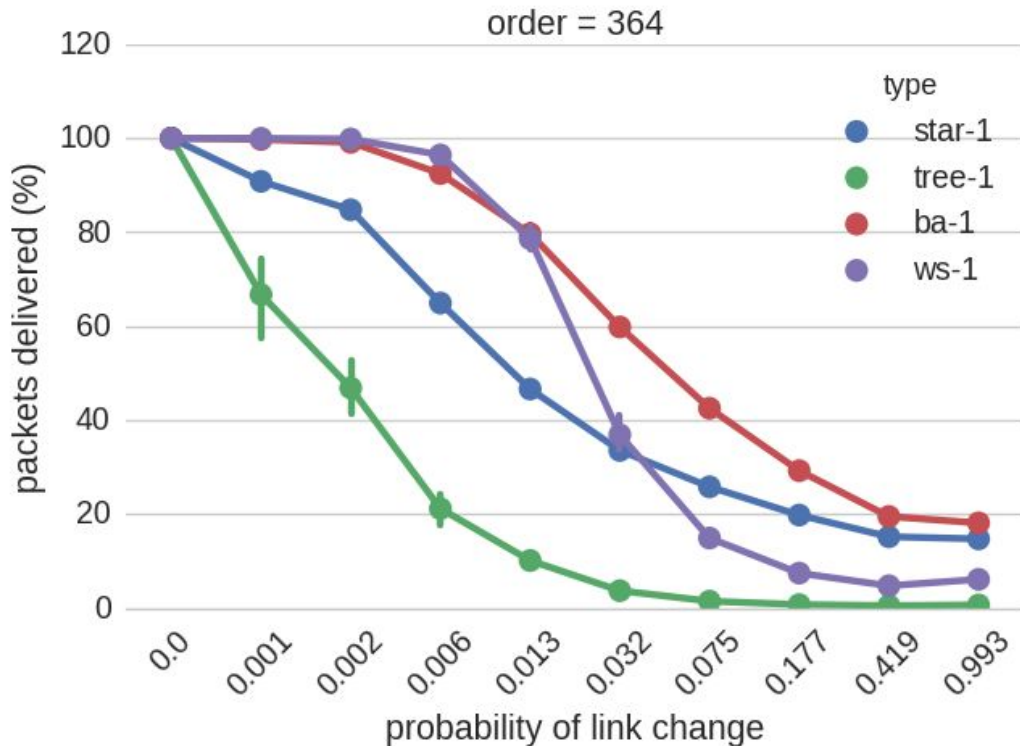
$$\frac{\textit{NO ROUTE dropped packets}}{\textit{packets transmitted}} \times 100 [\%]$$

Size effect



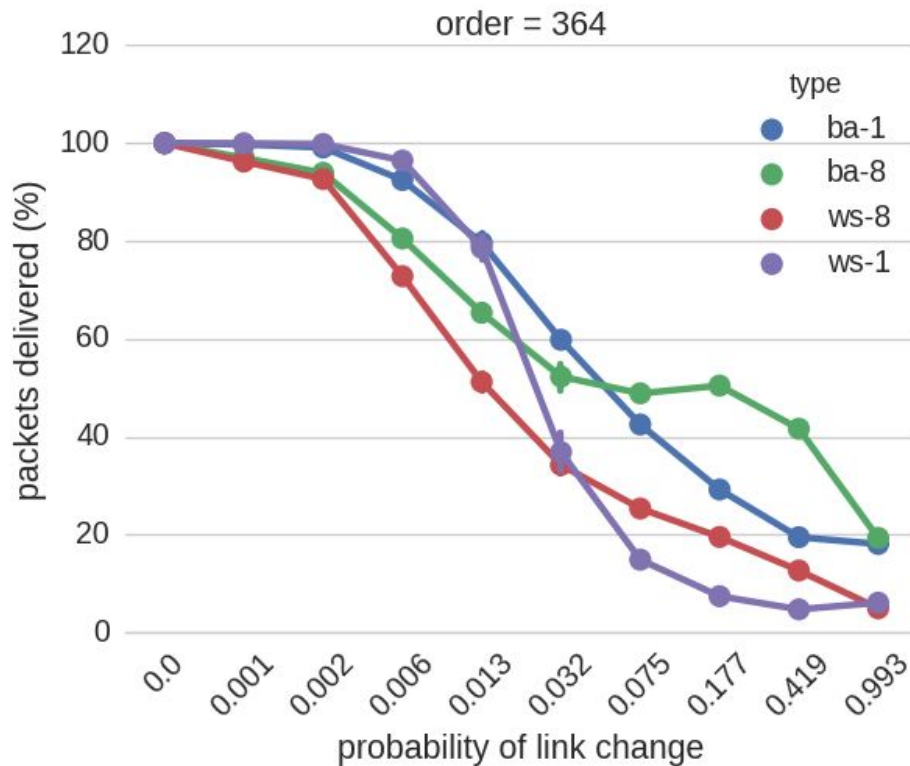
Topology **type** effect

- BA does significantly better for $p > 0.013$



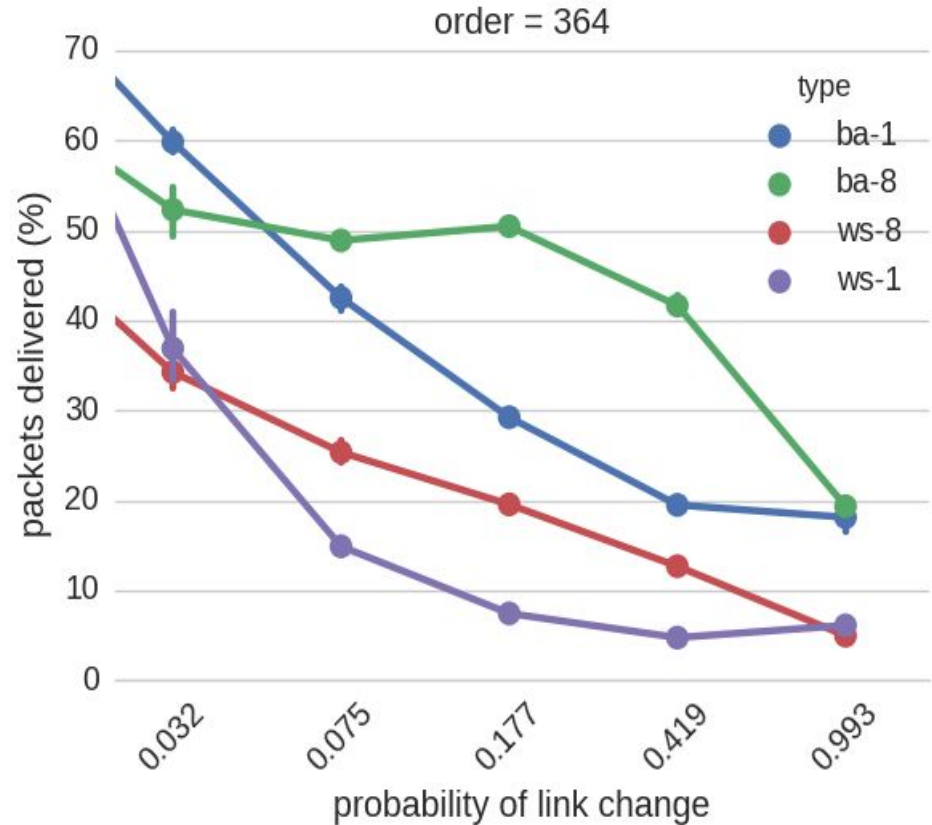
Partitioning effect

- For low link state change prob. ($p < 0.75$) it decreased the performance
- But, when the links started to become lossy...



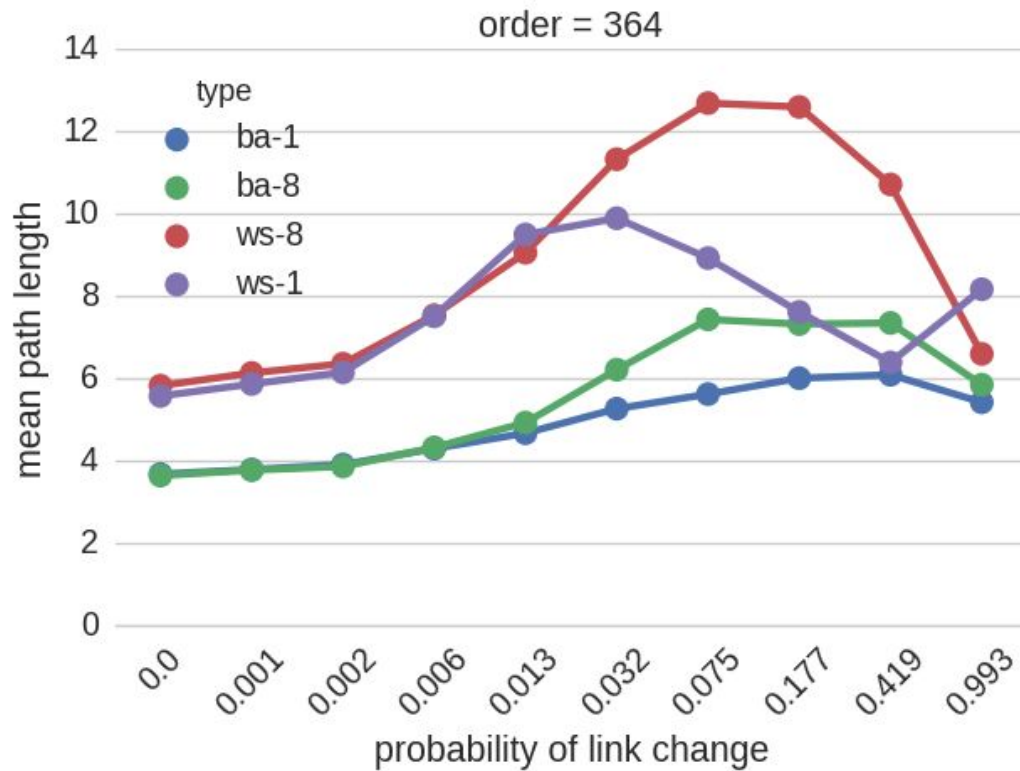
Partitioning effect (prob. ≥ 0.075)

- It improved data delivery for both WS and BA ($p \geq 0.075$)



Partitioning effect

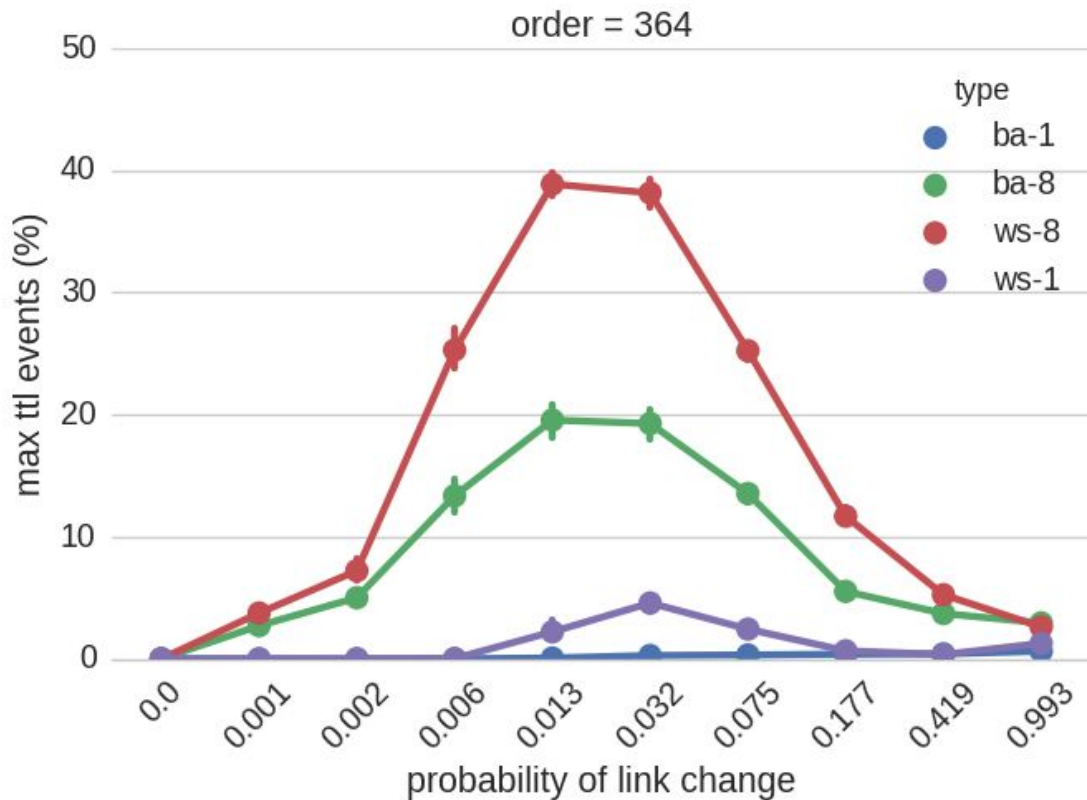
- Partitioning extended path length, $p > 0.013$



“TTL exceeded”

events

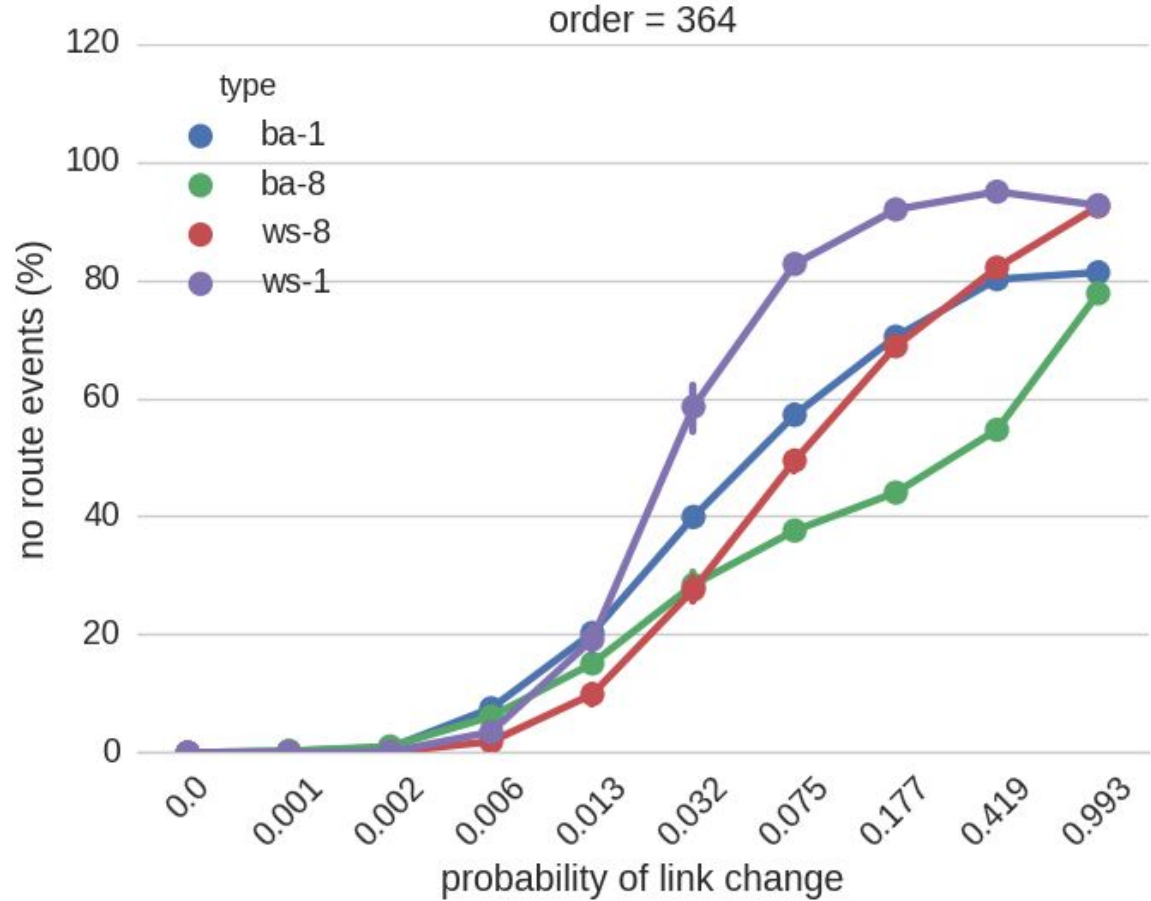
- Partitioning amplified loop occurrence
- Peak looping occurs near $p=0.013, 0.032$



“NO ROUTE”

events

- Less missing routes for partitioned nets



Conclusions

- **Size** increase caused lower packet delivery (more links, more failures)
- BA model **type** appears to be the most resilient (high number of edges, low diameter)
- Topology **partitioning** increased data delivery for “noisy” networks ($p > 0.075$)