

# Classical Cryptography

## Monoalphabetic cryptanalysis

Karst Koymans

Informatics Institute

University of Amsterdam

(version 22.4, 2023/02/13 13:41:58 UTC)

Tuesday, February 14, 2023

## 1 Statistical Cryptanalysis

- Frequencies
- The index of coincidence:  $\phi$ - and  $\chi$ -tests

## 2 Example

## 3 Countermeasures against statistical cryptanalysis

- Homophones
- Polyalphabetic substitutions

# Outline

- 1 Statistical Cryptanalysis
  - Frequencies
  - The index of coincidence:  $\phi$ - and  $\chi$ -tests
- 2 Example
- 3 Countermeasures against statistical cryptanalysis
  - Homophones
  - Polyalphabetic substitutions

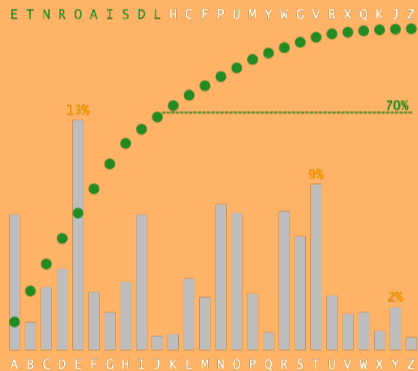
# Outline

- 1 Statistical Cryptanalysis
  - Frequencies
    - The index of coincidence:  $\phi$ - and  $\chi$ -tests
- 2 Example
- 3 Countermeasures against statistical cryptanalysis
  - Homophones
  - Polyalphabetic substitutions

# Letter frequencies

- A simple method to attack monoalphabetic ciphers
  - **Letter frequency analysis**
- Some letters occur more (or less) than others
  - This is (somewhat) language dependent

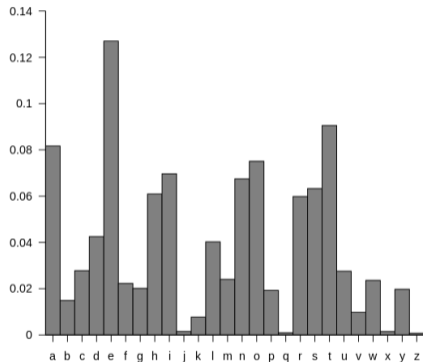
# Letter frequency diagram



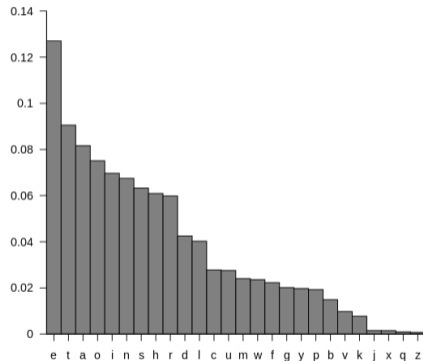
Source: Slides Hans van der Meer

Unknown language or text source

# English letter frequency



Ordered by alphabet



Ordered by frequency

Source: [https://en.wikipedia.org/wiki/Letter\\_frequency](https://en.wikipedia.org/wiki/Letter_frequency)

# Outline

- 1 Statistical Cryptanalysis
  - Frequencies
  - The index of coincidence:  $\phi$ - and  $\chi$ -tests
- 2 Example
- 3 Countermeasures against statistical cryptanalysis
  - Homophones
  - Polyalphabetic substitutions



# William Friedman



William Friedman

# The index of coincidence (IoC)

- Introduced by **William Friedman**
- Probability that two letters chosen randomly from a text, based on an alphabet of  $n$  letters, are the same
- Given probabilities of occurrence  $p_0, \dots, p_{n-1}$  for the  $n$  letters
  - $\text{IoC} = \sum_{i=0}^{n-1} p_i^2$
- For text with a (uniformly) random frequency distribution  
this reduces theoretically (obviously) to  $1/n$  ( $\approx 0.038$  for  $n = 26$ )
- For an English text (with the English frequency distribution)  
this amounts to  $\approx 0.066$ , or  $\approx 1/15$ , found by doing experiments

# The $\phi$ -test

- The IoC clearly distinguishes English text from random text
- Friedman observed that the IoC is **invariant under monoalphabetic substitution**
- Using the IoC to check for monoalphabeticity is called the  $\phi$ -test
- For an unknown ciphertext of length  $M > 1$  this test calculates
  - $\text{IoC} = \sum_{t=A}^Z f_t(f_t - 1) / M(M - 1)$
  - Here  $f_t$  is the number of occurrences of the letter  $t$
  - For small texts the  $-1$  is used to avoid counting identity as equality
    - Hence letters that occur only once don't contribute to the IoC

# Breaking Caesar (by hand and automatically)

- Brute force 26 keys and see if you get plaintext (we did this before)
- Match (visually) the frequency distribution of the cryptogram to standard English by shifting the frequency graph
- To automate this the  $\phi$ -test doesn't help, use the  $\chi$ -test instead
  - The  $\chi$ -test is also called cross-product sum
  - Consider two texts  $f$  and  $g$  of length  $M$  and  $N$  and calculate  $\chi = \sum_{t=A}^Z f_t g_t / MN$
  - Find the highest  $\chi$  value after comparing the shifted frequency diagram of the cryptogram with that of normal English text

# Breaking general monoalphabetic substitutions

- First use the  $\phi$ -test to check for monoalphabeticity
- Order the ciphertext letter distribution by frequency and try to match this with the standard English letter distribution or whatever language you may suspect is being used
- Look at digraph – or even trigraph – frequencies
- Look at beginning and ending of words (each has a different frequency distribution)
- Check vowels versus consonants and other letter patterns
- Look at keywords for alphabet construction
- Try to find cribs

# Outline

- 1 Statistical Cryptanalysis
  - Frequencies
  - The index of coincidence:  $\phi$ - and  $\chi$ -tests
- 2 Example
- 3 Countermeasures against statistical cryptanalysis
  - Homophones
  - Polyalphabetic substitutions

## Math of Secrets: 2.2 monoalphabet

QBVDL WXTEQ GXOKT NGZJQ GKXST RQLYR  
XJYGJ NALRX OTQLS LRKJQ FJYGJ NGXLK  
QLYUZ GJSXQ GXSLQ XNQXL VXKOJ DVJNN  
BTKJZ BKPXU LYUNZ XLQXU JYQGX NTYQG  
XKXQJ KXULK QJNQN LQBYL OLKKX SJYQG  
XNGLU XRSBN XOFUL YDSXU GJNSX DNVTY  
RGXUG JNLEE SXLYU ESLYY XUQGX NSLTD  
GQXKB AVBKX JYYBR XYQNQ GXKXZ LNYBS  
LRPBA VLQXK JLSOB FNGLE EXYXU LSBYD  
XWXKF SJQQS XZGJS XQGXF RLVXQ BMXXK  
OTQKX VLJYX UQBZG JQXZL NG

# Exercise 1

## Exercise 1

- Count letters and make a table of frequencies
- Generate a frequency diagram, using a spreadsheet
- Calculate the Index of Coincidence
- Is it an additive cipher?
- Try to solve the cryptogram by assuming it is affine



# Outline

- 1 Statistical Cryptanalysis
  - Frequencies
  - The index of coincidence:  $\phi$ - and  $\chi$ -tests
- 2 Example
- 3 Countermeasures against statistical cryptanalysis
  - Homophones
  - Polyalphabetic substitutions

# Outline

- 1 Statistical Cryptanalysis
  - Frequencies
  - The index of coincidence:  $\phi$ - and  $\chi$ -tests
- 2 Example
- 3 Countermeasures against statistical cryptanalysis
  - **Homophones**
  - Polyalphabetic substitutions

# Homophones

- Homophones
  - A classic way to flatten frequency distributions
  - Introduce more than one ciphertext letter option for (some of) the plaintext letters
    - Especially for plaintext letters with high frequency
    - Needs a larger ciphertext alphabet
  - This is an example where the encryption function may be randomised (to a small extent)
  - The Zodiac Killer used homophones in both Z408 and Z340
    - But could have done a better job in randomisation

## Math of Secrets: 2.2 homophones

IW\*CI W@G\*L &H&L( ASN\*A E)U&V \$CNPC  
SIW\*E DDSA@ LTCIH !(A#C V%EIW \*!#HA  
\*IW@N TAEHR \$CI(C JTS!C SHDS# SIW@S  
DVW@R G\$HH\* SIW\*W )JH@( CUGDC IDUIW  
\*&AIP GWTUA TLS\$L CIW\*D IWTG! #HATW  
TRG\$H H\*SQT U\$G\*I W@S)D GHWTR APBDG  
\*S%EI W@WDB @HIG@ IRWWX H&CV+ XHWVG  
\*LLXI WW#HE G)VG@ HHI#A AEGTH @CIAN  
W\*L!H Q%I!L )DAAN R)BTI B)K#C VXC#I  
HDGQX ILXIW IW@VA \*&B!C SIWTH E\*\*S\$  
UA(VW I

## Exercise 2

### Exercise 2

- Count symbols and make a table of frequencies
- Generate a frequency diagram, using a spreadsheet
- Calculate the Index of Coincidence for all symbols
- Calculate the Index of Coincidence for only the letters
- Is it a monoalphabetic cipher?
- Identify homophones and solve the cryptogram

# Outline

- 1 Statistical Cryptanalysis
  - Frequencies
  - The index of coincidence:  $\phi$ - and  $\chi$ -tests
- 2 Example
- 3 Countermeasures against statistical cryptanalysis
  - Homophones
  - Polyalphabetic substitutions

# Polyalphabetic substitutions

## Definition (polygraphic)

A **polygraphic substitution** is the replacement of groups of letters by other groups of letters according to one big substitution table

## Definition (polyliteral)

A **polyliteral substitution** is the replacement of single letters by groups of letters according to one big substitution table

## Definition (polyalphabetic)

A **polyalphabetic substitution** is the replacement of single letters by other letters by using a **varying ciphertext alphabet** for encrypting each plaintext letter